Exercise Sheet 1 COMS10017 Algorithms 2020/2021

Reminder: $\log n$ denotes the binary logarithm, i.e., $\log n = \log_2 n$.

1 O-notation: Part I

Give formal proofs of the following statements using the definition of Big-O from the lecture (i.e., identify positive constants c, n_0 for which the definition holds):

- 1. $5\sqrt{n} \in O(n)$. 2. $n^2 + 10n + 8 \in O(\frac{1}{2}n^2)$. 3. $n^3 + n^2 + n = O(n^3)$. 4. $10 \in O(1)$.
- 5. $\sum_{i=1}^{n} i \in O(4n^2)$.

2 Racetrack Principle

Use the racetrack principle to prove the following statement:

 $n \leq e^n$ holds for every $n \geq 1$.

3 O-notation: Part II

Give formal proofs of the following statements using the definition of Big-O from the lecture.

- 1. $f \in O(h_1), g \in O(h_2)$ then $f \cdot g \in O(h_1 \cdot h_2)$.
- 2. $2^n \in O(n!)$.
- 3. $2^{\sqrt{\log n}} \in O(n)$.

4 Fast Peak Finding

Consider the following variant of FAST-PEAK-FINDING where the " \geq " sign in the condition in instruction 4 is replaced by a "<" sign:

- 1. if A is of length 1 then return 0
- 2. if A is of length 2 then compare A[0] and A[1] and return position of larger element
- 3. if $A[\lfloor n/2 \rfloor]$ is a peak then return $\lfloor n/2 \rfloor$
- 4. Otherwise, if $A[\lfloor n/2 \rfloor 1] < A[\lfloor n/2 \rfloor]$ then return FAST-PEAK-FINDING $(A[0, \lfloor n/2 \rfloor - 1])$
- 5. else
 - return $\lfloor n/2 \rfloor + 1 + \text{FAST-PEAK-FINDING}(A[\lfloor n/2 \rfloor + 1, n 1])$

Give an input array of length 8 on which this algorithm fails.

5 Optional and Difficult

Exercises in this section are intentionally more difficult and are there to challenge yourself.

5.1 Advanced Racetrack Principle

Use the racetrack principle and determine a value n_0 such that

$$\frac{2}{\log n} \leq \frac{1}{\log \log n} \text{ holds for every } n \geq n_0 \ .$$

Hint: Transform the inequality and eliminate the log-function from one side of the inequality before applying the racetrack principle. If needed, apply the racetrack principle twice! Recall that $(\log n)' = \frac{1}{n \ln(2)}$. The inequality $\ln(2) \ge 1/2$ may also be useful.

5.2 Finding Two Peaks

We are given an integer array A of length n that has exactly two peaks. The goal is to find both peaks. We could do this as follows: Simply go through the array with a loop and check every array element. This strategy has a runtime of O(n) (requires $c \cdot n$ array accesses, for some constant c). Is there a faster algorithm for this problem (e.g. similar to FAST-PEAK-FINDING)? If yes, give such an algorithm. If no, justify why there is no such algorithm.