

Exercise Sheet 3

COMS10017 Algorithms 2020/2021

Reminder: $\log n$ denotes the binary logarithm, i.e., $\log n = \log_2 n$.

1 Warm up: Proof by Induction

Consider the following sequence: $s_1 = 1$, $s_2 = 2$, $s_3 = 3$, and $s_n = s_{n-1} + s_{n-2} + s_{n-3}$, for every $n \geq 4$. Prove that the following holds:

$$s_n \leq 2^n .$$

2 Loop Invariant

Prove that the stated invariant holds throughout the execution of the loop (using the Initialization, Maintenance, Termination approach discussed in the lectures):

Algorithm 1 Algorithm \mathcal{A}

Require: Array A of length n ($n \geq 2$)

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1:  $S \leftarrow A[0] - A[1]$ 
2: for  $i \leftarrow 1 \dots n - 2$  do
3:    $S \leftarrow S + A[i] - A[i + 1]$ 
4: end for
5: return  $S$ 
```

Invariant:

At the beginning of iteration i , $S = A[0] - A[i]$ holds.

What does the algorithm compute?

3 Insertionsort

What is the runtime (in Θ -notation) of Insertionsort when executed on the following arrays of lengths n :

- 1, 2, 3, 4, \dots , $n - 1$, n
- n , $n - 1$, $n - 2$, \dots , 2, 1
- The array A such that $A[i] = 1$ if $i \in \{1, 2, 4, 8, 16, \dots\}$ (i.e., when i is a power of two) and $A[i] = i$ otherwise.

4 Mergesort

The Mergesort algorithm uses the MERGE operation, which assumes that the left and the right halves of an array A are already sorted, and merges these two halves so that A is sorted afterwards. The runtime of this operation is $O(n)$.

Suppose that we replaced the MERGE operation in our Mergesort algorithm with a less efficient implementation that runs in time $O(n^2)$ (instead of $O(n)$). What is the runtime of our modified Mergesort algorithm?

5 Runtime Analysis

Algorithm 2

Require: Integer $n \geq 2$

$x \leftarrow 0$

$i \leftarrow n$

while $i \geq 2$ **do**

$j \leftarrow \lceil n^{1/4} \rceil \cdot i$

while $j \geq i$ **do**

$x \leftarrow x + 1$

$j \leftarrow j - 10$

end while

$i \leftarrow \lfloor i/\sqrt{n} \rfloor$

end while

return x

Determine the runtime of Algorithm 3 in Θ -notation.