

Heapsort

COMS10018 - Algorithms

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Sorting Algorithms seen so far

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- Insertionsort: $O(n^2)$ in worst case, in place, stable
- Mergesort: $O(n \log n)$ in worst case, NOT in place, stable

Heapsort (best of the two)

- $O(n \log n)$ in worst case, in place, **NOT** stable
- Uses a *heap data structure* (a heap is special tree)

Data Structures

- *Data storage format that allows for efficient access and modification*
- Building block of many efficient algorithms
- For example, an array is a data structure

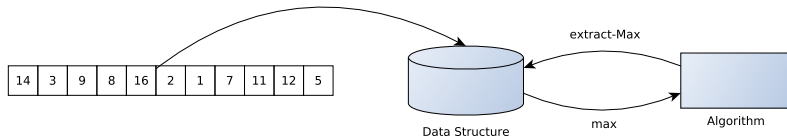
Priority Queues

Priority Queue:

Data structure that allows the following operations:

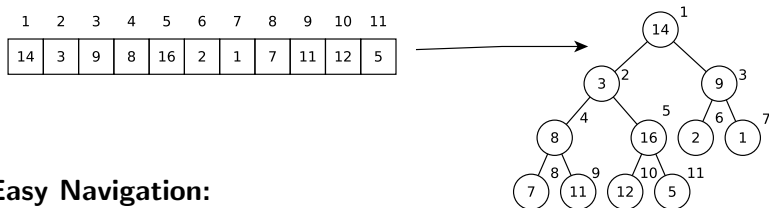
- `Create(.)`: Create data structure given a set of data items
- `Extract-Max(.)`: Remove the maximum element from the data structure and return it
- *others...*

Sorting using a Priority Queue



From Array to Tree

Interpretation of an Array as a Complete Binary Tree

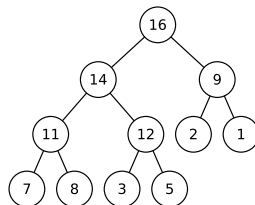
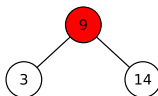
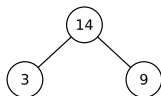


Easy Navigation:

- Parent of i : $\lfloor i/2 \rfloor$
- Left/Right Child of i : $2i$ and $2i + 1$

The Heap Property

Key of nodes larger than keys of their children



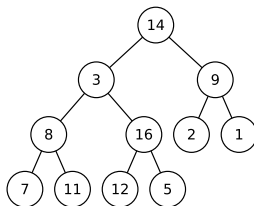
Heap Property \rightarrow Maximum at root
Important for Extract-Max(.)

The Heapify Operation

Constructing a Heap: Create-Heap(.)

Given a binary tree, transform it into one that fulfills the Heap Property

- 1 Traverse tree with regards to right-to-left array ordering
- 2 If node does not fulfill Heap Property: **Heapify()**

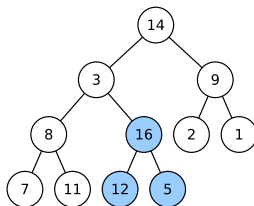


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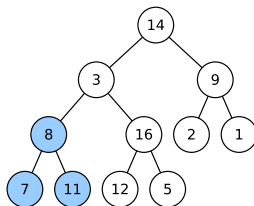


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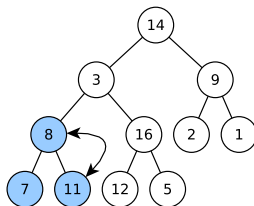


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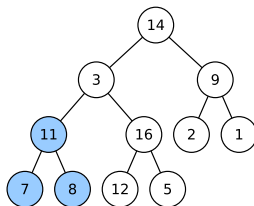


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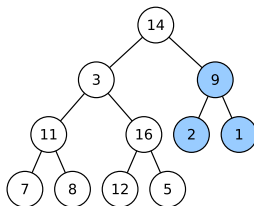


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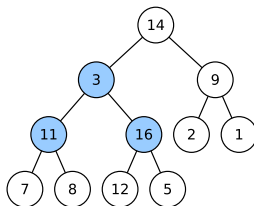


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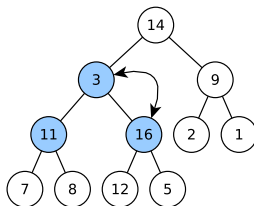


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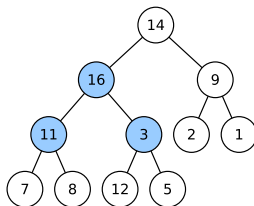


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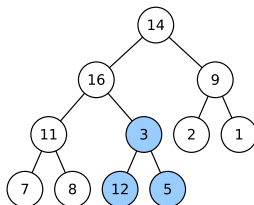


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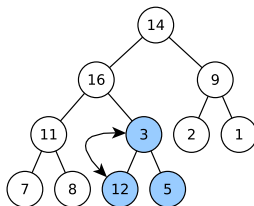


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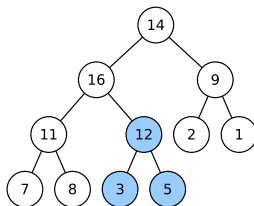


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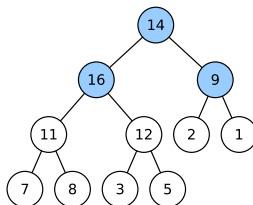


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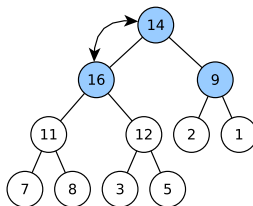


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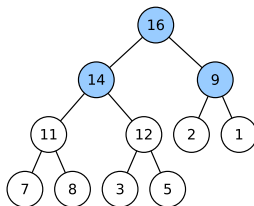


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Runtime of Heapify()

Heapify()

Let p be the key of a node and let c_1, c_2 be the keys of its children

- Let $c = \max\{c_1, c_2\}$
- If $c > p$ then exchange nodes with keys p and c
- call **Heapify()** recursively at node with key p

Runtime:

- Exchanging nodes requires time $O(1)$
- The number of recursive calls is bounded by the height of the tree, i.e., $O(\log n)$
- Runtime of **Heapify**: $O(\log n)$.

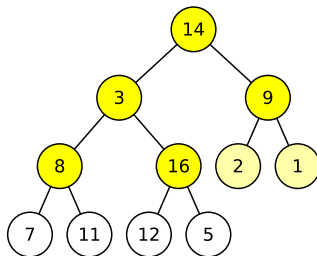
Constructing a Heap: Create-Heap(.) Runtime $O(n \log n)$

More Precise Analysis of the Heap Construction Step

- $\text{Heapify}(x)$: $O(\text{depth of subtree rooted at } x) = O(\log n)$
- **Observe:** Most nodes close to the “bottom” in a complete binary tree

Analysis:

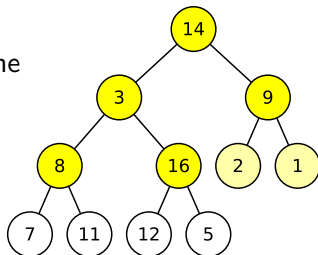
- Let i be the largest integer such that $n' := 2^i - 1$ and $n' < n$
- There are at most n' internal nodes (candidates for $\text{Heapify}()$)
- These nodes are contained in a perfect binary tree
- This tree has i levels



Improved Analysis of Heap Construction

Analysis

We sum over all relevant levels, count the number of nodes per level, and multiply with the depth of their subtrees:



$$\begin{aligned}\text{Runtime} &= \sum_{j=1}^i \# \text{ nodes at level } (i - j + 1) \cdot \text{depth of subtree} \cdot O(1) \\ &= O(1) \sum_{j=1}^i 2^{i-j} \cdot j = O(1) \cdot 2^i \cdot \sum_{j=1}^i \frac{j}{2^j} \\ &= O(2^i) = O(n') = O(n) ,\end{aligned}$$

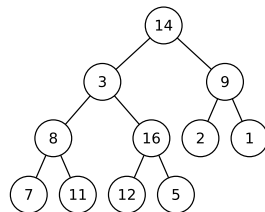
using $\sum_{j=1}^i \frac{j}{2^j} = O(1)$ (see trick from linear/binary search lecture).

The Complete Algorithm

Putting Everything Together

14	3	9	8	16	2	1	7	11	12	5
----	---	---	---	----	---	---	---	----	----	---

- 1 Create-Heap()
- 2 Repeat n times:
 - 1 Swap root with last element
 - 2 Decrease size of heap by 1
 - 3 Heapify(root)

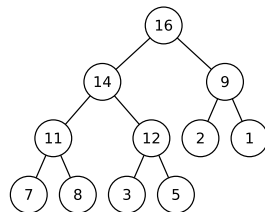


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16	14	9	11	12	2	1	7	8	3	5
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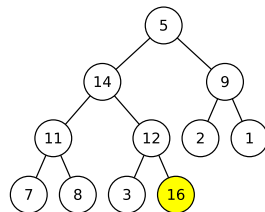


The Complete Algorithm

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5	14	9	11	12	2	1	7	8	3	16
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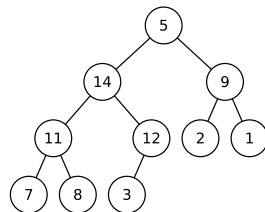


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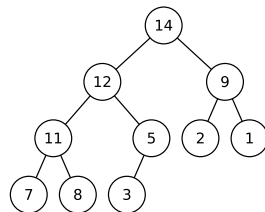


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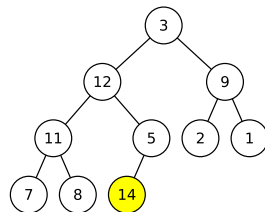


The Complete Algorithm

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3	12	9	11	5	2	1	7	8	14	16
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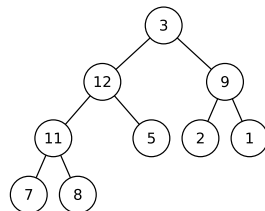


The Complete Algorithm

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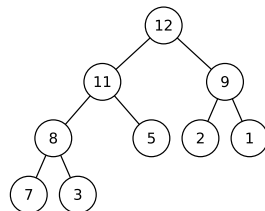


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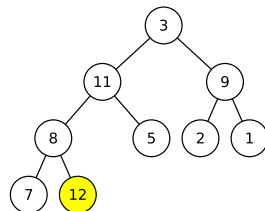


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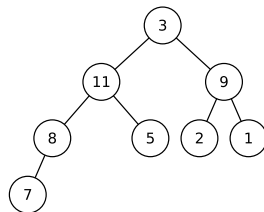


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...

Putting Everything Together

1	2	3	5	7	8	9	11	12	14	16
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- ❶ Create-Heap()
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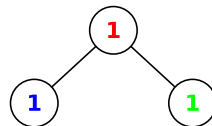
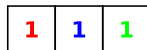
- ❶ Create-Heap() $O(n)$
- ❷ Repeat n times:
 - ❶ Swap root with last element $O(1)$
 - ❷ Decrease size of heap by 1 $O(1)$
 - ❸ Heapify(root) $O(\log n)$

Runtime: $O(n \log n)$

Heapsort is Not Stable

Example:

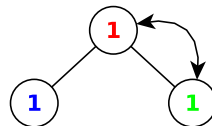
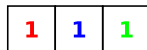
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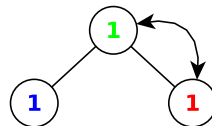
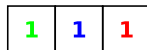
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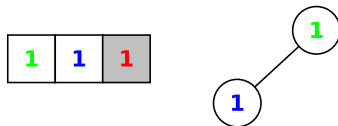
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1 is moved from left to the right past 1 and 1

Heap-sort not stable